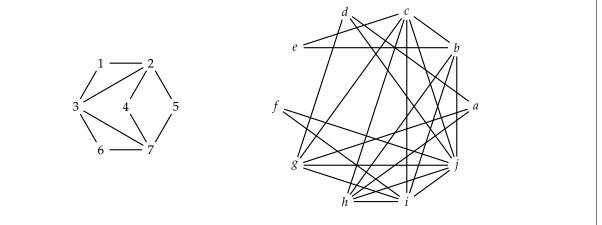
## Model 1: Independent sets

**Definition 1.** An *independent set* in an undirected graph G = (V, E) is a subset of vertices  $S \subseteq V$  such that no two vertices in *S* are adjacent.

**Definition 2.** A *vertex cover* in an undirected graph G = (V, E) is a subset of vertices  $C \subseteq V$  such that every edge  $e \in E$  has at least one endpoint in (is "covered by") *C*.



- 1 Which of the following are independent sets?
- (a) {1,2}
- (b) {1,5}
- (c)  $\{c, a\}$
- (d)  $\{e, a, i, g\}$
- (e)  $\{7\}$
- (f) Ø
- 2 For each graph, list at least three other examples of independent sets.
- 3 Given an arbitrary graph *G*, does *G* always have at least one independent set? Why or why not?

Think of edges as hallways in an art museum and *C* as the set of locations where we are going to put some guards. Then an independent set means no two guards can see each other; a vertex cover means every hallway is watched by at least one guard.

- 4 Intuitively, which is harder: to find big independent sets, or small ones? Why?
- 5 Based on the previous observation, an interesting question to ask

about a giver	ı graph G is to fi	nd the
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- 6 Try to answer your interesting question for the given example graphs (but don't spend more than a few minutes). How sure are you about your answer?
- 7 Describe a brute-force algorithm to answer this question. What is its big- $\Theta$  running time in terms of |V| and |E|?
- 8 Guess the running time (in terms of |V| and |E|) of the fastest known algorithm to solve this problem. (You do not have to come up with an algorithm; just guess how fast you think this problem can be solved.)
- 9 Which of the following are vertex covers?
- (a)  $\{3, 4, 5, 6, 7\}$
- (b)  $\{2, 3, 4, 6, 7\}$
- (c)  $\{b, d, e, f, g, h, i, j\}$
- (d)  $\{b, c, d, f, h, j\}$
- (e)  $\{1, 2, 3, 4, 5, 6\}$
- (f)  $\{1, 2, 3, 4, 5, 6, 7\}$

10 For each graph, list at least three other examples of vertex covers.



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- 11 Given an arbitrary graph *G*, does *G* always have at least one vertex cover? Why or why not?
- 12 Intuitively, which is harder: to find small vertex covers, or big ones? Why?
- 13 Based on the previous observation, an interesting question to ask

about a given graph *G* is to find the \_\_\_\_\_.

- 14 Try to answer your interesting question for the given example graphs. How sure are you about your answer?
- 15 Describe a brute-force algorithm to answer this question. What is its big- $\Theta$  running time in terms of |V| and |E|?
- 16 Compare your answers to questions 1 and 9. What do you notice?

Make a conjecture based on your observations in the previous section:

**Theorem 3.** Let G = (V, E) be an undirected graph, and  $S \subseteq V$  a subset of

its vertices. Then S is an independent set if and only if \_\_\_\_\_.



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Let's prove it! This requires proving "both directions" of the claim. For the first direction, a skeleton proof is provided. For the reverse direction, you must write the proof from scratch.

*Proof.*  $(\Longrightarrow)$  Let *S* be an independent set. We must show

So pick an arbitrary edge $e = (u, v) \in$	E E;	
by definition we must show that at least one of <i>u</i> or <i>v</i>	/	
that is, at least one of $u$ or $v$ is not Since $S$ is an independent set and $u$ and $v$ are connected by an		
edge, <i>u</i> and <i>v</i> can't both	/	
and therefore		
Now, fill in the proof for the "other" direction!. $(\Leftarrow)$		Write down what you get to assume and what you are trying to prove, and expand definitions.

