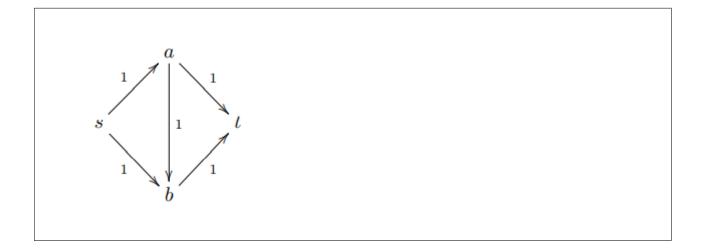
## Algorithms: Max Flow

- 1 Consider the graph in Model 1 and the algorithm below. The algorithm claims to find in a graph *V*, *E* the maximum flow  $f(e) \forall e \in E$  through the graph. On line 3, it arbitrarily chooses an unsaturated path and adds flow to it. What is an ordering of choices that enables this algorithm to successfully find a maximum flow?
  - 1. Initialize  $f(e) \leftarrow 0$  for all  $e \in E$
  - 2. repeat

3. Find an unsaturated path P from s to t

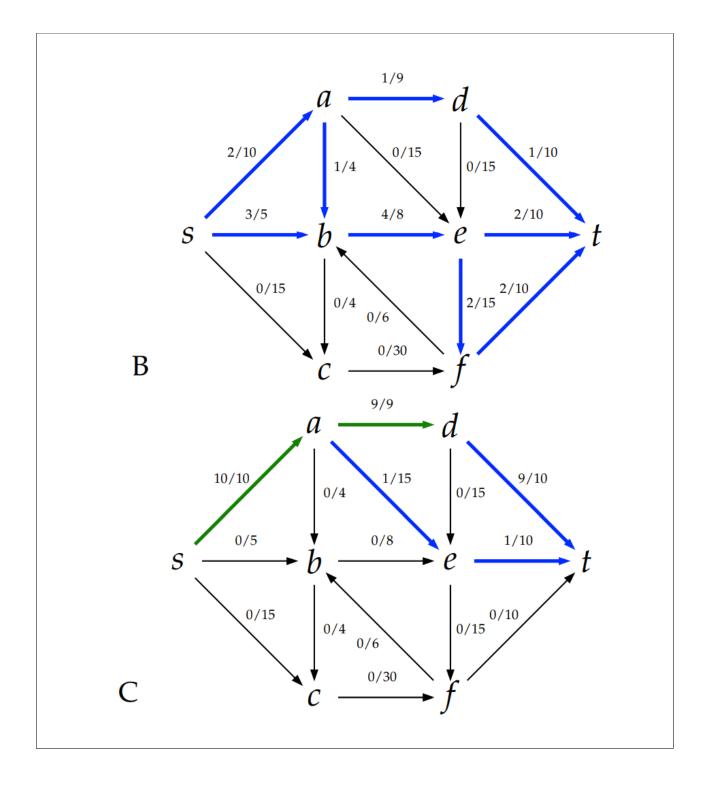
- 4.  $a \leftarrow \text{minimum excess capacity } c(e) f(e) \text{ among all}$ edges  $e \in P$
- 5.  $f(e) \leftarrow f(e) + a$  for each edge  $e \in P$
- 6. until no more unsaturated  $s \leftarrow t$  paths
- 2 What is an ordering of choices that prevents this algorithm from finding a maximum flow?

Model 1: Graph



**Learning objective**: Students will derive an efficient algorithm for finding the flow of maximum value in a graph.

Model 2: Flow Graphs



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**Definition 1.** Given a network G = (V, E) and a flow f on G, we define the *residual graph*  $G_f$  as  $G_f = (V, E_f)$ , with

$$E_f = \{e \in E \mid c(e) - f(e) > 0\} \cup \{e^R \mid e \in E, f(e) > 0\}$$

(where  $e^R$  denotes the reverse of edge e, i.e., if e = (u, v), then  $e^R = (v, u)$ ) and we define the capacities of edges in  $E_f$  by

- $c_f(e) = c(e) f(e)$ , and
- $c_f(e^R) = f(e).$
- 3 Based on the given definition of a residual graph, draw a residual graph for each of the flow graphs depicted in Model 2.



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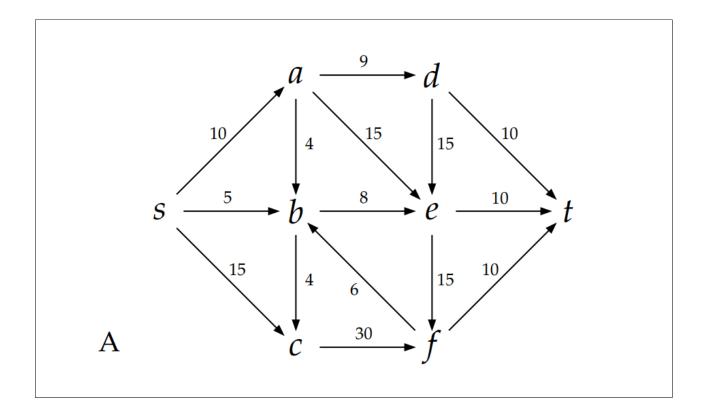
## 4 Consider the algorithm below.

- 1. Initialize  $f(e) \leftarrow 0$  for all  $e \in E$
- 2. repeat
- 3.  $\alpha \leftarrow \min\{c_f(e) \mid e \in P\}$
- 4.  $f(e) \leftarrow f(e) + \alpha$  for each  $e \in P$  such that  $e \in E$
- 5.  $f(e) \leftarrow f(e) \alpha$  for each  $e \in P$  such that  $e^R \in E$
- 6. until no paths exist from s to t in  $G_f$

This algorithm, the Ford-Fulkerson algorithm, uses the residual graph  $G_f$  instead of the original graph to find *augmenting paths*. We execute lines 3, 4 and 5 because we have found an augmenting path in the residual graph.

Use the Ford-Fulkerson algorithm to find a flow on the graph in Model 3. Select paths in any order you like.

## Model 3: Original Graph





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5 Is the flow you computed in the previous question the maximum possible flow? Why or why not?



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