

Stable Matching Problem :: Definitions

Problem statement. Given a set H of n hospitals, a set D of n doctors, and their respective preferences, can we construct a *perfect matching* of doctors to hospitals that is *stable*—that is, no doctor-hospital pair has an incentive to break their match.

Perfect Matching: A matching M is perfect if it assigns each doctor to a single hospital and each hospital to a single doctor.

Stable matching: A matching M is *stable* if there is no pair $(d, h) \in H \times D$ where both:

1. d prefers h to its current match in M , and
2. h prefers d to its current match in M .

PROPOSE-REJECT Algorithm by Gale Shapley

Let's examine the Gale-Shapley propose-reject algorithm:

Algorithm 1 PROPOSE-REJECT - finds a stable matching

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1: Initialize each doctor  $d$  and hospital  $h$  as FREE
2: while there is a free doctor who hasn't proposed to every hospital do
3:   Choose a free doctor  $d$ 
4:    $h \leftarrow$  first hospital on  $d$ 's list to whom  $d$  has not yet proposed
5:   if  $h$  is FREE then
6:      $d$  and  $h$  are MATCHED
7:   else if  $h$  prefers  $d$  to its current match  $d'$  then
8:      $d$  and  $h$  are MATCHED and  $d'$  is FREE
9:   else
10:     $h$  rejects  $d$  and  $d$  remains FREE
11:   end if
12: end while
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