

Type Inference

CSCI 334
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Type Checking

- Dynamic
 - + Prevents Errors
 - run-time overhead
 - catch errors only at run time
 - coverage: only catch errors on paths executed
 - + more flexible
- Static
 - + Prevents Errors
 - + no run-time cost
 - + catch errors at compile time
 - + coverage: catches errors on all possible paths
 - restricts flexibility

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Limitations of Static Type Checking

- Types restrict expressiveness:
 - Lisp: '(1 "cow" 4.5)
 - ML: [1,2,3] : int list
 - zip_any in ML
- Decidability

```
fun f() {  
  while (big-test) {  
    if (big-test-2) {  
      return 3;  
    }  
  }  
  return 3 + true;  
}  
  
fun h(int x) {  
  int a[] = new int[10];  
  int y = (x*x+sqrt(x)*2)/g(x);  
  a[y] = 3;  
}
```

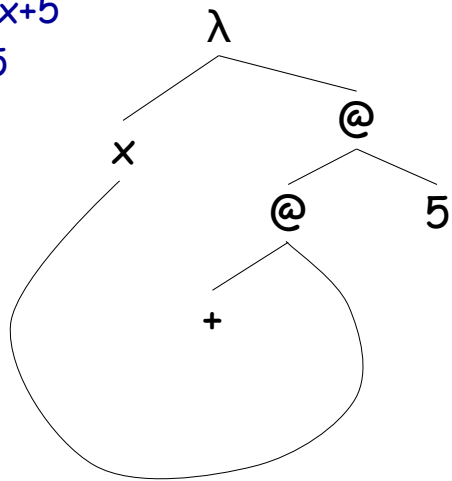
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Checking vs. Inference

- Type Checking
 - int f(int x) { return x + 1; }
 - int g(int y) { return f(y) * 2; }
- Type Inference
 - fun f(x) = x + 1;
 - fun g(y) = f(y) * 2;

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fun g(x) = x+5
 g = λx.x+5



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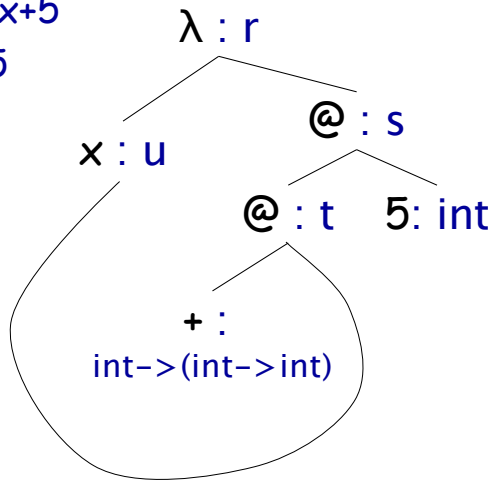
Hindley-Milner Algorithm

• **Goal:** Annotate each node in tree with a type

1. Annotate nodes with
 - actual type for constants, ops, known names
 - unique type variable for all others
2. Generate constraints over type variables using structure of tree.
3. Solve constraints via unification

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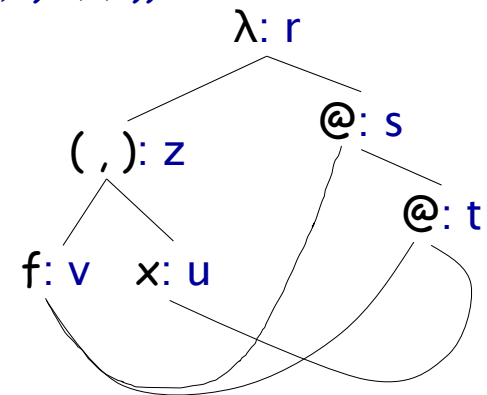
fun g(x) = x+5
 g = λx.x+5



Step 1

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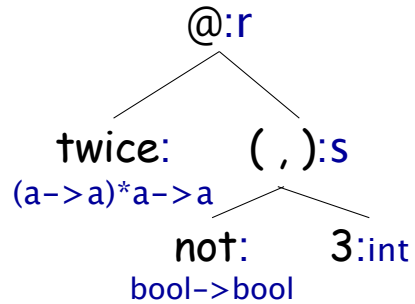
fun twice(f,x) = f(f(x))
 twice = λ(f,x).f(f(x))



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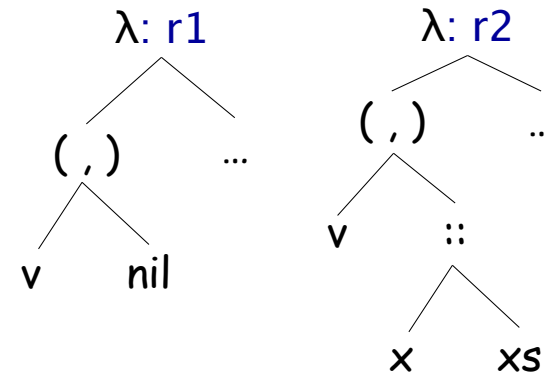
not: bool -> bool
twice(not, 3);

s = (bool -> bool) * int
(a -> a) * a -> a = s -> r
(a -> a) * a = s
a = r
(a -> a) * a = (bool -> bool) * int
(a -> a) = (bool -> bool)
a = int
a = bool



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fun search(v,nil) = ...
| search(v,x::xs) = ...



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Java Autoboxing

```
class Vector<E> {
    E elementData[];
    int elementCount;

    void add(E o) {
        elementData[elementCount++] = o;
    }
}

Vector<Integer> v = new Vector <Integer>();
v.add(3);
...
println(v.get(0) * 2);
```

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Java Reality: Boxing/Unboxing

```
class Vector {
    Object elementData[];
    int elementCount;

    void add(Object o) {
        elementData[elementCount++] = o;
    }
}

Vector v = new Vector();
v.add(new Integer(3));
...
println(((Integer)v.get(0)).intValue()*2);
```

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C++ Templates

```
template <typename T> void sort(T d[], int n) {
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < n - i - 1; j++) {
            if (d[j] < d[j+1]) {
                T temp = d[j];
                d[j] = d[j+1];
                d[j+1] = temp;
            }
        }
    }
}

int a[100];
char b[100];
double c[100];
sort(a,100);
sort(b,100);
sort(c,100);
```

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add for ints

```
template <typename T> T add(T x, T y)
{ return x+y; }

add(1,2);

add<int>:
    pushl    %ebp
    movl    %esp, %ebp
    movl    12(%ebp), %eax
    addl    8(%ebp), %eax
    leave
    ret
```

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add for reals

```
add(1.3,2.3);

add<double>:
    pushl    %ebp
    movl    %esp, %ebp
    movsd   %xmm0, -8(%rbp)
    movsd   %xmm1, -16(%rbp)
    movsd   -8(%rbp), %xmm0
    addsd   -16(%rbp), %xmm0
    popq    %rbp
    retq
```

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Code Specialization

```
int f(int g, int r, int x, int y) {
    return g*g*g*x + y/(r*r);
}

int main() {
    for (int i = 0; i < 50000000; i++) {
        int a = f(100, 100, i, 30);
        int b = f(10, 90, 1, 3);
        int c = f(20, 20, i, 3);
        ...
    }
}
```

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Code Specialization

```
template <int G, int R> int f(int x, int y) {  
    return G*G*G*x + y/(R*R);  
}  
  
int main() {  
    for (int i = 0; i < 50000000; i++) {  
        int a = f<100,100>(i,30);  
        int b = f<10,90>(1,3);  
        int c = f<20,20>(i,3);  
        ...  
    }  
}
```

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Code Specialization

```
template <int G, int R> int f(int x, int y) {  
    return G*G*G*x + y/(R*R);  
}  
  
f0(int x, int y) {  
    return 1000000*x + y/10000;  
}  
  
int main() {  
    for (int i = 0; i < 50000000; i++) {  
        int a = f0(i,30);  
        int b = f1(1,3);  
        int c = f2(i,3);  
        ...  
    }  
}  
  
f2(int x, int y) {  
    return 8000*x + y/400;  
}
```

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